

Cooperation and communication dynamics

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PSE

What do we model?

Dynamic interactions

- An agent's decisions at some point in the game influence other agent's behaviors in the future.
- This part of the course is mostly concerned with long-term interactions.
- Drawn on the theory of repeated games.

Natural applications

- Interaction between firms,
- consumer and producer,
- repeated contracting,
- social norms, customs etc.

What are we after and how?

Cooperation

- Under what circumstances can repetition lead to **cooperation**?
- More generally, how does repetition affect **sustainable outcomes**?
- What **social norms** do we expect in repeated interactions?

Information dynamics

- What is the best way to use private information?
 - Using information reveals it, making it less valuable in the future.
 - Should one use information or conceal it?
 - Is all private information revealed in the long-run?
- Is it possible to create anticipations for a certain behavior?
 - Can a central bank, creditor, producer, or chef build a reputation?
 - How does reputation building influence interactions?
 - Are reputations always good?
 - Do reputations survive in the long-run?

Methodology

Rules of the game

- What do the players initially know?
- What do they observe as the game is played?
- Define strategies in the repeated interaction
- Define preferences (payoffs) in the repeated interaction

Solving and interpreting

- Use equilibrium concepts: Nash equilibria and refinements,
- Draw predictions on player's behavior and on equilibrium payoffs
- Interpret. Are we satisfied with the predictions of the theory?
- If not, is the model or solution concept not appropriate?

Roadmap

- 1 A few examples
- 2 Modeling
- 3 Finitely repeated games

A first example

	<i>H</i>	<i>U</i>
<i>G</i>	1, 1	-1, -1
<i>S</i>	2, 0	-1, -1

- What are the Nash equilibria of this game?
- Consider the game played twice. After the first interaction, actions in the first interaction are observed.
 - Draw the game in extensive form
 - What is a strategy for player 1? For player 2? How many of them are there?
 - Do (or don't) represent the game in normal form
 - What are the Nash equilibria?
 - How does the presence of a dominated strategy affect the Nash equilibria of the repeated game?
 - What are the subgames?
 - What are the SPNE of the repeated game?

A second example

	<i>H</i>	<i>U</i>	<i>B</i>
<i>G</i>	1, 1	-1, -1	-5, -5
<i>S</i>	2, 0	-1, -1	-5, -5
<i>B</i>	-5, -5	-5, -5	-4, -4

- What are the Nash equilibria of this game?
- Consider the game played twice. After the first interaction, actions in the first interaction are observed.
 - Do (or don't) represent the game in extensive form
 - How many strategies for player i ?
 - Do (or don't) represent the game in normal form
 - What are the SPNE of the repeated game?
 - How does the presence of the B strategy influence the SPNE?

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Stage game

Stage game

- set I of players, $i \in I$,
- set A^i of actions for player i ,
- $A = \prod_i A^i$, $A^{-i} = \prod_{j \neq i} A^j$,
- payoff function for player i : $g^i: A \rightarrow \mathbb{R}$,
- vector payoff function g , $g(a) = (g^i(a))_i$,

Mixed strategies

- $S^i = \Delta(A^i)$, $S = \prod_i S^i$, $S^{-i} = \prod_{j \neq i} S^j$,
- extended payoff function: $g^i: S \rightarrow \mathbb{R}$, $g^i(s) = (g^i(s))_i$,

$$g^i(s) = \mathbf{E}_s g^i(a)$$

- Nash equilibrium: $\tilde{s} \in S$ such that for every i , $s^i \in S^i$

$$g^i(s^i, \tilde{s}^{-i}) \leq g^i(\tilde{s})$$

Information, histories, pure strategies

Histories

- The set of histories of length t is $H_t = A^t$.
- By convention, $H_0 = \{\emptyset\}$ (or $\{*\}$, or anything you want).
- An element $h_t = (a_1, \dots, a_t) \in H_t$ describes the information available to all the players prior to playing in stage $t + 1$.

Pure strategies

- Player i 's behavior at stage t is described by a map $f_t^i: H_{t-1} \rightarrow A^i$
- A pure strategy for i in the (infinitely) repeated game is $f^i = (f_t^i)_{t \geq 1}$
- Equivalently, it is a mapping $f^i: \cup_{t \geq 0} H_t \rightarrow A^i$
- Let Σ^i be the set of pure strategies for player i

The “right” notion of strategies

What is a randomized strategy ?

- A **mixed strategy** for player i is a randomization over pure strategies, element of $\Delta(\Sigma^i)$.
- A **behavioral strategy** is a mapping from $\cup_{t \geq 0} H_t$ to $\Delta(A^i)$ describing the randomization at each stage.
- More generally, a randomized strategy is a randomization over behavioral strategies.
- So what is the “right” notion?

Kuhn's Theorem

In a game with **perfect recall**, for every mixed strategy of player i , there exists an equivalent behavioral strategy of player i , and vice versa. Equivalent means, no matter the behavior of other players, the probability distributions induced over histories are the same.

Repeated games **without** perfect recall are also interesting! (bounded rationality), but beyond the scope of this course.

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Induced histories, preferences

Histories

A profile f of p.s. induces a play $h_\infty = (a_1, \dots, a_t, \dots)$ in the r.g.

- $a_1 = f(\emptyset) = (f^i(\emptyset))_i$, $h_1 = (a_1)$,
- $a_2 = f(h_1)$, $h_2 = (a_1, a_2)$
- $a_t = f(h_{t-1})$, $h_t = (a_1, \dots, a_t)$

Payoffs

- G_n game played n times,

$$\gamma_n = \frac{1}{n} \sum_{t=1}^n g(a_t)$$

- E_n set of Nash Equilibrium (vector) payoffs of G_n
- What can we say about E_n . for a fixed value of n , for n large?

Example, Prisoner's Dilemma

	<i>C</i>	<i>D</i>
<i>C</i>	3, 3	-1, 4
<i>D</i>	4, -1	0, 0

What are E_1 , E_2 , E_n ?

A first general (negative) result

Defending a payoff

Player i can **defend** a payoff x^i in G if, for every s^{-i} in S^{-i} , there exists $s^i \in S^i$ such that

$$g^i(s^{-i}, s^i) \geq x^i$$

Anti-Folk Theorem

Assume that G has a unique Nash Equilibrium payoff x , and that for every i , player i can defend x^i , then for every n ,

$$E_n = \{x\}.$$

Proof

Let f be a NE of G_n in (wlog, pure) strategies, inducing a sequence of actions a_1, \dots, a_n . We show by induction on k that $\forall k, g(a_{n-k}) = x$.

- true for $k = 0$
- $k - 1 \rightarrow k$, assume a_{n-k} is not a NE of G , find a profitable deviation for a player in G_n .

Subgames and subgame perfection

Subgames and continuation strategies

- To each history h_τ , $\tau \leq n - 1$ corresponds a subgame of G_n
- A strategy f^i for player i induces a strategy $f_{|h_\tau}^i$ in the repeated game starting **after** h_τ given by

$$f_{|h_\tau}^i(h_t) = f^i(h_\tau \cdot h_t)$$

where $h_\tau \cdot h_t$ is the history h_τ followed by h_t .

Subgame perfect Nash equilibria

- f is a SPNE of G_n if, for every history h_τ , $\tau \leq n - 1$, the profile of continuation strategies $f_{|h_\tau} = (f_{|h_\tau}^i)_i$ is a Nash equilibrium of $G_{n-\tau}$.
- E'_n denotes the set of SPNE payoffs of G_n .

Back to an example and a general result

	H	U
G	1, 1	-1, -1
S	2, 0	-1, -1

The game G_2 has a Nash equilibrium which is not a SPNE. In all SPNE of G_2 , players play (S, H) at every stage, and after every history.

Can we generalize the result to G_n ?

Compare with the SPNE of the two stage repetition of

	H	U	B
G	1, 1	-1, -1	-5, -5
S	2, 0	-1, -1	-5, -5
B	-5, -5	-5, -5	-4, -4

Anti-Folk part deux

Assume that G has a unique Nash equilibrium, \tilde{s} . Then, G_n has a unique SPNE, and at this SPNE, \tilde{s} is played after every history.

Conclusion

- Repetition **can** lead to cooperation in a repeated game,
- “end of game” effect can lead to **unraveling** and prevent cooperation
- Repetition does not **necessarily** lead to cooperation
- How about infinite (never-ending) games?